

# Hierarchical Decision-Making in Population Games

Yu-Wen Chen<sup>1</sup>, Nuno C. Martins<sup>2</sup>, Murat Arcak<sup>1</sup>

<sup>1</sup>Department of Electrical Engineering and Computer Sciences, UC Berkeley

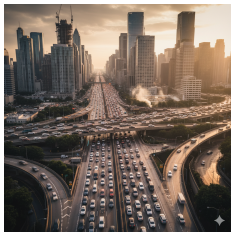
<sup>2</sup>Department of Electrical and Computer Engineering, University of Maryland

December 11, 2025



paper/slides

# Large-scale strategic multi-agent systems (MASs)



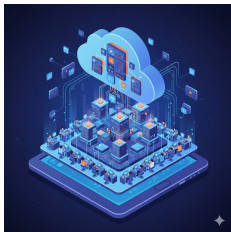
(a) transportation



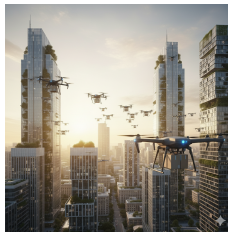
(b) smart grids



(c) markets



(d) resource allocations

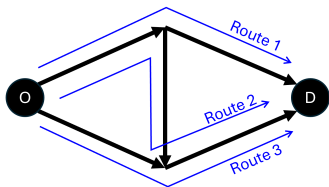


(e) Drone-as-a-Service

figures generated with Google Gemini

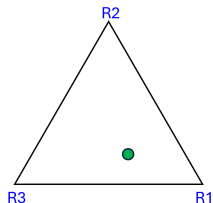
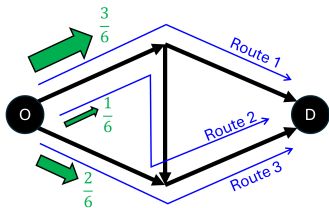
# Large-scale strategic MASs: population games

- a large number of homogeneous players
- a **strategy set**  $\mathcal{S} = \{1, \dots, n\}$



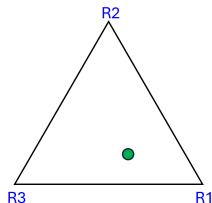
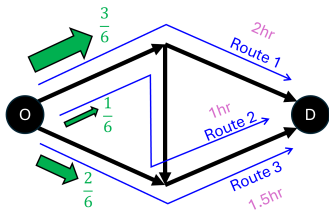
# Large-scale strategic MASs: population games

- a large number of homogeneous players
- a **strategy set**  $\mathcal{S} = \{1, \dots, n\}$
- a **social state**  $\mathbf{x} = (x_1, \dots, x_n) \in \Delta^n := \{x_i \geq 0, \sum_i x_i = 1\}$



# Large-scale strategic MASs: population games

- a large number of homogeneous players
- a **strategy set**  $\mathcal{S} = \{1, \dots, n\}$
- a **social state**  $\mathbf{x} = (x_1, \dots, x_n) \in \Delta^n := \{x_i \geq 0, \sum_i x_i = 1\}$
- a **payoff function**  $F : \Delta^n \rightarrow \mathbb{R}^n$  maps  $\mathbf{x}$  to payoffs  $[\dots, F_i(\mathbf{x}), \dots]^T$

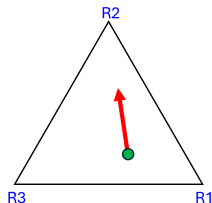
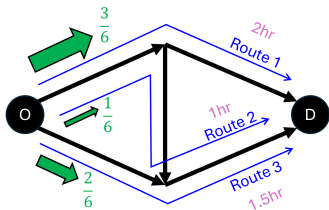


# Large-scale strategic MASs: population games

- a large number of homogeneous players
- a **strategy set**  $\mathcal{S} = \{1, \dots, n\}$
- a **social state**  $\mathbf{x} = (x_1, \dots, x_n) \in \Delta^n := \{x_i \geq 0, \sum_i x_i = 1\}$
- a **payoff function**  $F : \Delta^n \rightarrow \mathbb{R}^n$  maps  $\mathbf{x}$  to payoffs  $[\dots, F_i(\mathbf{x}), \dots]^T$
- an **incentive-oriented revision protocol**<sup>1</sup>  $\mathcal{T} : \Delta^n \times \mathbb{R}^n \rightarrow \mathbb{R}_+^{n \times n}$ , where  $\mathcal{T}_{ij}$  is the switching rate from strategy  $i$  to  $j$ , inducing mean dynamics

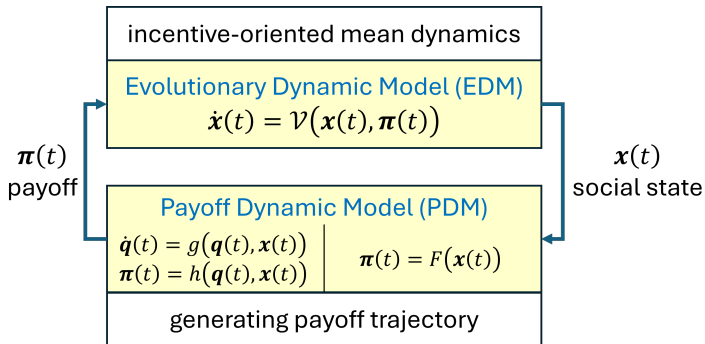
$$\dot{x}_i = \sum_{j=1}^n x_j \mathcal{T}_{ji}(\mathbf{x}, F(\mathbf{x})) - x_i \mathcal{T}_{ij}(\mathbf{x}, F(\mathbf{x})), \quad i \in \mathcal{S},$$

which we call the **evolutionary dynamics model (EDM)**



<sup>1</sup>Best Response, Pairwise Comparison, Imitation, Excess Payoff

# Main studies in population games: Convergence

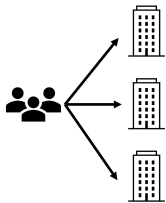


- with EDMs, what properties of PDMs lead to convergence?
  - not asking agents to follow some designed policies
  - not only about finding equilibrium, but about convergence of dynamics

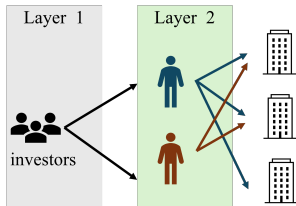
# Overview of our paper: hierarchical population games

- Motivation:

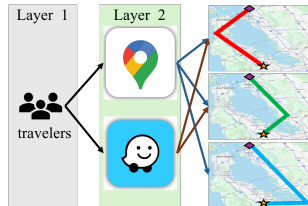
Most population games assume that agents make direct decisions.



(a) direct framework



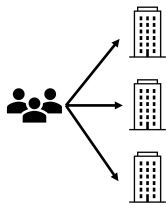
(b) get info & make a decision



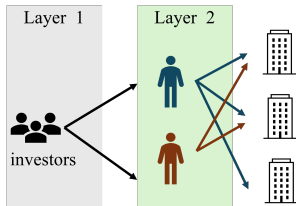
(c) get info & make a decision

# Overview of our paper: hierarchical population games

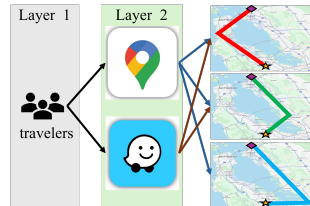
- Motivation:  
Most population games assume that agents make direct decisions.
- Contribution:
  - We propose a hierarchical decision-making framework for population games.
  - We prove convergence of system dynamics under some payoff structures.



(a) direct framework

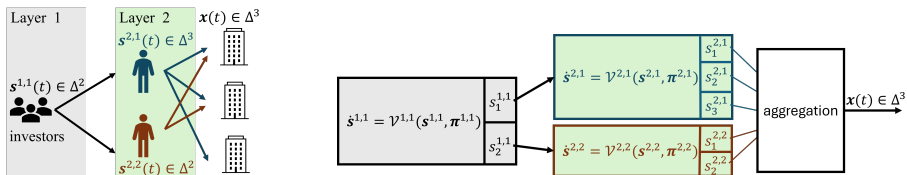


(b) get info & make a decision



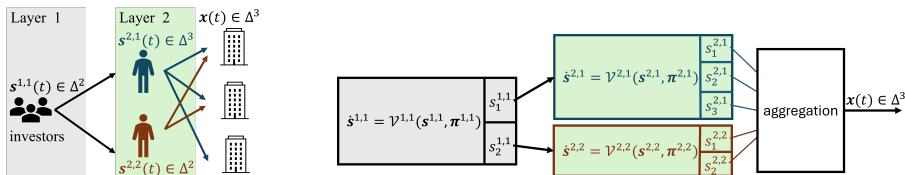
(c) get info & make a decision

# Proposed basic hierarchical structure



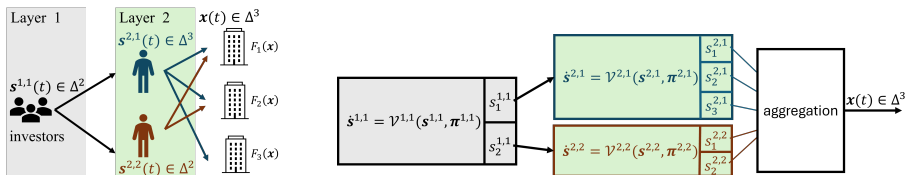
- each layer  $i = 1, \dots, L$  has  $n^i$  groups indexed by  $j = 1, \dots, n^i$
- $s^{i,j} \in \Delta^{d^{i,j}}$  is the distribution of  $(i, j)$ -group over its  $d^{i,j}$  strategies

# Proposed basic hierarchical structure



- each layer  $i = 1, \dots, L$  has  $n^i$  groups indexed by  $j = 1, \dots, n^i$
- $s^{i,j} \in \Delta^{d^{i,j}}$  is the distribution of  $(i, j)$ -group over its  $d^{i,j}$  strategies
- social state  $x(t) = \begin{bmatrix} s_1^{1,1}(t)s_1^{2,1}(t) + s_2^{1,1}(t)s_1^{2,2}(t) \\ s_1^{1,1}(t)s_2^{2,1}(t) + s_2^{1,1}(t)s_2^{2,2}(t) \\ s_1^{1,1}(t)s_3^{2,1}(t) \end{bmatrix}$ : percentage to targets

# Proposed basic hierarchical structure



- each layer  $i = 1, \dots, L$  has  $n^i$  groups indexed by  $j = 1, \dots, n^i$
- $s^{i,j} \in \Delta^{d^{i,j}}$  is the distribution of  $(i, j)$ -group over its  $d^{i,j}$  strategies

- social state  $\mathbf{x}(t) = \begin{bmatrix} s_1^{1,1}(t)s_1^{2,1}(t) + s_2^{1,1}(t)s_1^{2,2}(t) \\ s_1^{1,1}(t)s_2^{2,1}(t) + s_2^{1,1}(t)s_2^{2,2}(t) \\ s_1^{1,1}(t)s_3^{2,1}(t) \end{bmatrix}$ : percentage to targets

- payoff  $\pi^{i,j}(t)$  is (backwardly) determined by the average performance

$$\pi^{2,1}(t) = F(\mathbf{x}(t)),$$

$F(\cdot)$ : social payoff function

$$\pi^{2,2}(t) = [F_1(\mathbf{x}(t)), F_2(\mathbf{x}(t))]^T,$$

$$F(\mathbf{x}) = [F_1(\mathbf{x}), F_2(\mathbf{x}), F_3(\mathbf{x})]^T$$

$$\pi^{1,1}(t) = \left[ \sum_{l=1}^3 s_l^{2,1}(t)F_l(\mathbf{x}(t)), \sum_{l=1}^2 s_l^{2,2}(t)F_l(\mathbf{x}(t)) \right]^T$$

# Strategy distribution of interest (allowable set) $K^{i,j}$

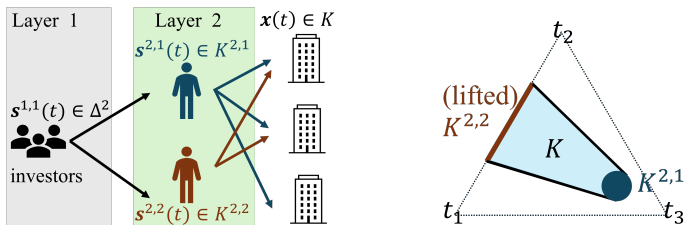
- We consider general cases:  $s^{i,j} \in K^{i,j} \subseteq \Delta^{d^{i,j}}$ , e.g.,
  - Manager 1 group** is restricted to the distribution  $s_*^{2,1} = [0.2, 0.2, 0.6]^T$  over the three investment targets but allows variations up to  $\varepsilon = 0.1$ ,

$$K^{2,1} = \{x \in \Delta^3 : \|x - s_*^{2,1}\|_2 \leq \varepsilon\}.$$

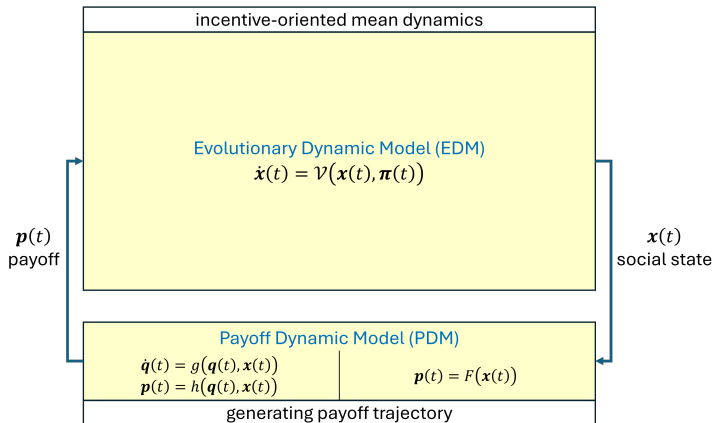
- Manager 2 group** needs to invest at least 1/3 in both investment targets,

$$K^{2,2} = \left\{ x \in \Delta^2 : x_1 \geq \frac{1}{3}, x_2 \geq \frac{1}{3} \right\}.$$

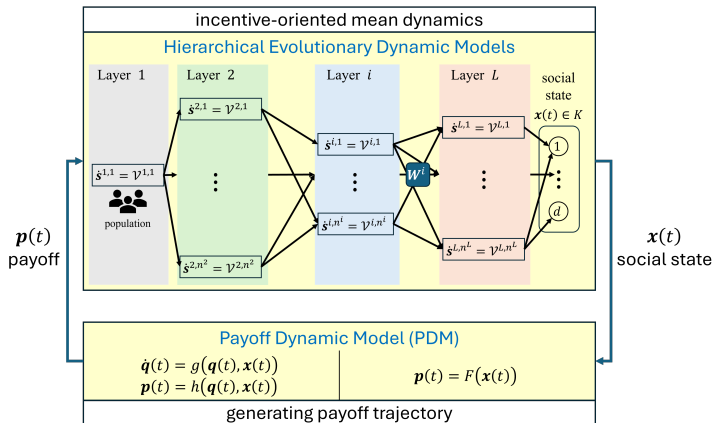
$\Rightarrow$  social state  $x(t)$  is confined within the admissible set  $K$ .



# Research focus on hierarchical population games

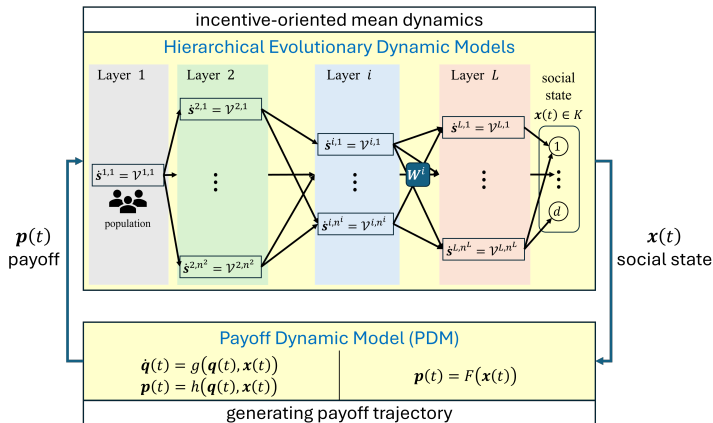


# Research focus on hierarchical population games



- properties on PDM s.t.  $x(t)$  evolves within  $K$  and converges
  - in classical case: agents react directly to the payoff function  $F$   
 $\Rightarrow$  rest points correspond to the set of Nash equilibria  $NE(F)$

# Research focus on hierarchical population games



- properties on PDM s.t.  $x(t)$  evolves within  $K$  and converges
  - in classical case: agents react directly to the payoff function  $F$   
 $\Rightarrow$  rest points correspond to the set of Nash equilibria  $NE(F)$
  - in hierarchical case: rest points meaningful? related to  $NE_K(F)$ ?

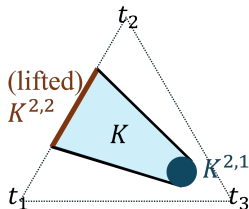
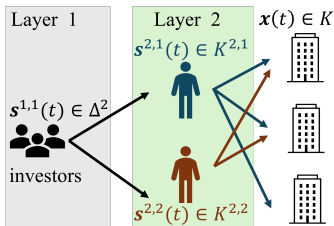
# Existence of $\text{NE}_K(F)$

## Theorem (Convexity of $K$ )

If  $K^{i,j}$ s are all convex, then  $K$  is convex.

⇒ By Kakutani's Fixed-Point Theorem:

The convexity of  $K$  guarantees the existence of  $\text{NE}_K(F)$ .

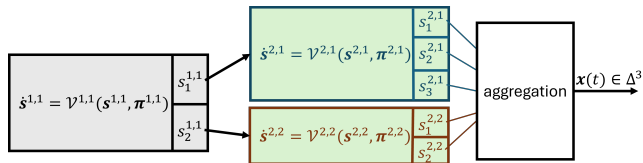


# Rest points are (refined) Nash equilibria

## Characterization at rest points (informal)

When each  $s^{i,j}$  attains a Nash equilibrium in its population game (w.r.t.  $\pi^{i,j}$ ), the corresponding social state  $x$  satisfies  $x \in \text{NE}_K(F)$ .

- all groups ( $s^{i,j}$ ) contribute to the social state ( $x$ )
- all groups are playing population games w/ different payoff functions

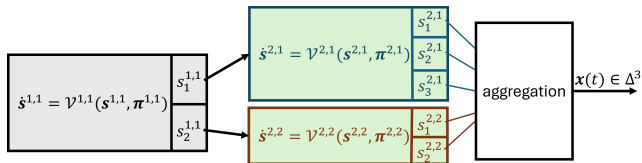


# Rest points are (refined) Nash equilibria

## Characterization at rest points (informal)

When each  $s^{i,j}$  attains a Nash equilibrium in its population game (w.r.t.  $\pi^{i,j}$ ), the corresponding social state  $x$  satisfies  $x \in \text{NE}_K(F)$ .

- all groups ( $s^{i,j}$ ) contribute to the social state ( $x$ )
  - all groups are playing population games w/ different payoff functions
- ⇒ the proposed **average payoff determination for  $\pi^{i,j}$**  renders that the decision groups, as players, participate in a (weighted) potential game



# Convergence results for hierarchical population games

## Definition (Nash Stationarity (NS) & Positive Correlation (PC))

An EDM is

- (1) NS w.r.t  $K$  if  $\mathcal{V}(s, \pi) = \mathbf{0} \iff s \in \beta_K(\pi) = \arg \max_{y \in K} y^T \pi$ .
- (2) PC if  $\mathcal{V}(s, \pi) \neq \mathbf{0} \implies \pi^T \mathcal{V}(s, \pi) > 0$ .

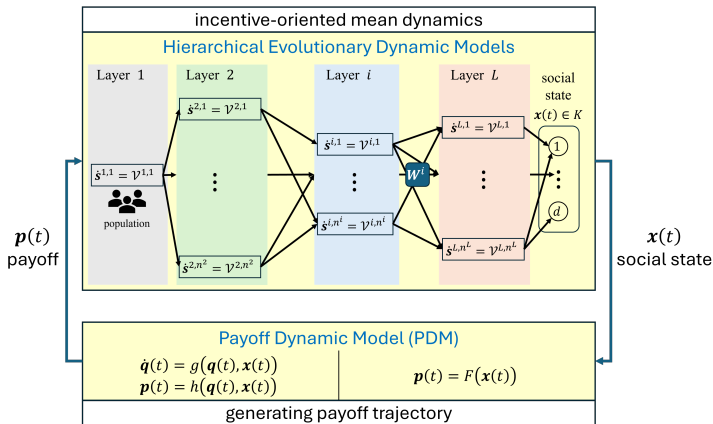
## Theorem (Convergence for potential game $F$ )

Let  $F$  be a potential game, and some regularities hold (Appendix). Suppose each  $\mathcal{V}^{i,j}$  is NS w.r.t.  $K^{i,j}$  and PC. Then,  $x(t)$  asymptotically approaches  $\text{NE}_K(F)$ .

## Theorem (Convergence for counterclockwise dissipativity PDM)

Let the payoff  $p(t)$  be given by a counterclockwise dissipativity PDM. Suppose that each  $\mathcal{V}^{i,j}$  is NS w.r.t.  $K^{i,j}$  and PC. Then,  $x(t) \rightarrow \beta_K(p(t))$ .

# Conclusions



paper/slides



website

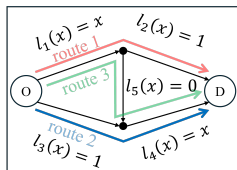
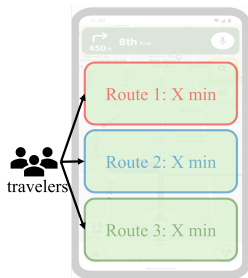
- Proposed a hierarchical decision-making structure for population games.
- Proved convergence under some payoff structures.

# Appendix

# Assumptions

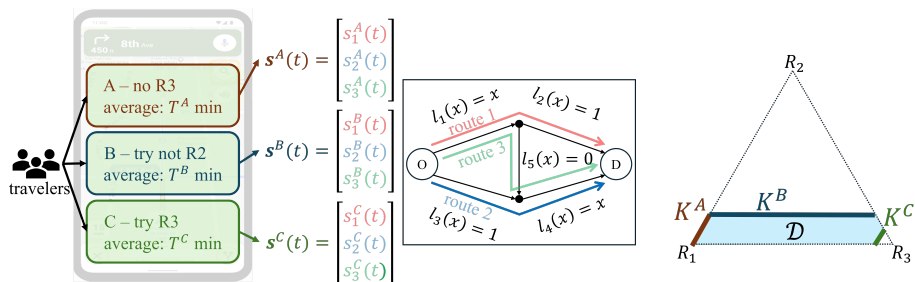
- 1  $K^{i,j}, i \in [L], j \in [n^i]$ , are compactly convex.
  - 2  $F$  is continuously differentiable.
  - 3 Each  $\mathcal{V}^{i,j}$  ensures the forward invariant of  $K^{i,j}$ .
  - 4  $\pi^T \mathcal{V}^{i,j}(s, \pi), i \in [L], j \in [n^i]$ , are Lipschitz continuous w.r.t.  $s$  and  $\pi$ .
- ⇒ When  $L = 1$  and  $K^{1,1} = \Delta^{d^{1,1}}$ , the assumptions reduce to the ones in classical population games.
- ⇒ Assumption 3 means that each  $(i, j)$ -group makes decisions consistently within its strategy distribution of interest  $K^{i,j}$ .

# As a design tool: partial information revealing



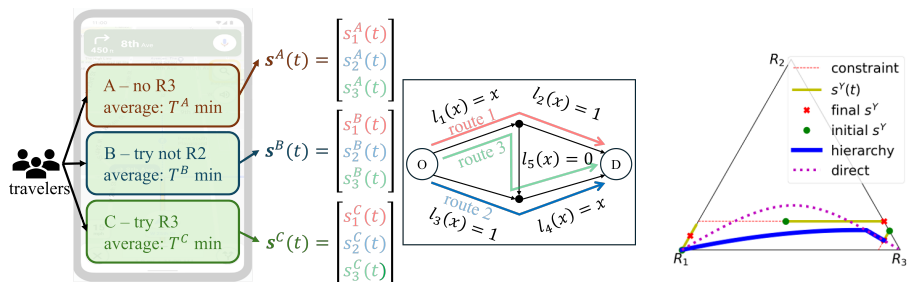
- If directly provide all information, then travelers face the Braess paradox.

# As a design tool: partial information revealing



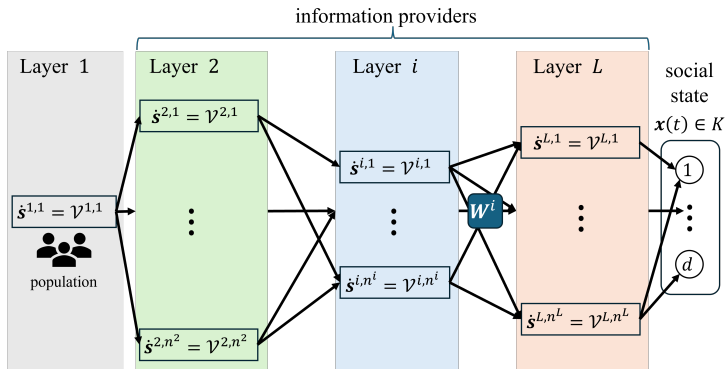
- If directly provide all information, then travelers face the Braess paradox.
- Instead, provide (verifiable) partial information, like:
  - each option plays a population game for users based on its known information
  - joining through option  $X$  undergoes average time  $T^X$ , for  $X = A, B, C$

# As a design tool: partial information revealing



- If directly provide all information, then travelers face the Braess paradox.
- Instead, provide (verifiable) partial information, like:
  - each option plays a population game for users based on its known information
  - joining through option  $X$  undergoes average time  $T^X$ , for  $X = A, B, C$

# Hierarchical population games bring more strategic topics



- agent level: revision protocol
- information providers: cooperation / competition